(1) Friday a starty BFS
Non basil venable ophnality condition
(3) [early readle fearible Condition] Basic variable -> non hour
(Gaussian Elimeter)
Assume that it is enterry. XBr to be deany which r.
$Y = B^{T}A \Leftrightarrow A = BY$
If you you you you was the entery consule the next schools fearable you have leave variable

$$\max f(\vec{x}) = \vec{C} \vec{x} \iff \lim_{x \to \infty} \lim_{x \to \infty} \int_{-\infty}^{\infty} |A\vec{x}| dx$$

$$f(\vec{x}) = f(\vec{x}_0) + \nabla f(\vec{x}_0)(\vec{x} - \vec{x}_0) + h.u.t.$$

$$(\Rightarrow) f(\vec{x}) = f(\vec{x}_0) + \nabla f(\vec{x}_0)(\vec{x} - \vec{x}_0) \times (\because h)_{\text{lin}}$$

$$f(\vec{x}) = f(\vec{x}_0) + \nabla^{\dagger}(\vec{x} - \vec{x}_0) \times (\because h)_{\text{lin}}$$

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$$f(\vec{x}_0) = f(\vec{x}_0) + \nabla^{\dagger}(\vec{x}_0 - \vec{x}_0) \times ((\vec{x}_0 - \vec{x}_0) \times ((\vec{x}_0 - \vec{x}_0) \times (\vec{x}_0 - \vec{x}_0) \times ((\vec{x}_0 - \vec{x}_0) \times (\vec{x}_0 - \vec{$$

$$A = \begin{bmatrix} B \mid R \end{bmatrix} \qquad \stackrel{\text{Convent}}{\times} = \begin{bmatrix} \stackrel{\text{X}}{\times} B \\ \stackrel{\text{Y}}{\circ} \end{bmatrix}$$

$$= \sum_{i=1}^{m} C_{Bi} \left(-X_{Bi} \frac{y_{ij}}{y_{rj}}\right) - C_{Bi} \cdot X_{Bi} \frac{y_{rj}}{y_{rj}} + C_{j} \cdot \frac{X_{Br}}{y_{rj}}$$

$$= \frac{1}{1-1} \left(C_{8i} \times K_{8r} \left(\frac{y_{5j}}{y_{rj}} \right) + C_{j} \left(\frac{x_{8r}}{y_{rj}} \right) \right)$$

$$=\frac{X_{Br}}{y_{G}}\left\{C_{j}-\sum_{i=1}^{m}C_{Bi}y_{ij}\right\}=\frac{X_{Br}}{y_{G}}\left\{C_{j}-Z_{i}\right\}$$

$$f(\vec{x}) = f(\vec{x}_0) + dod (n \times_j exp.)$$

$$= f(\vec{x}_0) + (x_0) + (x_0)$$

Since x is a feasible solution, Ax = b. Thus by (2.24),

$$\mathbf{b} = \sum_{j=1}^{n} x_{j} \mathbf{a}_{j} = \sum_{j=1}^{n} x_{j} (B \mathbf{y}_{j}) = \sum_{j=1}^{n} \left(\sum_{i=1}^{m} y_{ij} \mathbf{b}_{i} \right) x_{j} = \sum_{i=1}^{m} \left(\sum_{j=1}^{n} y_{ij} x_{j} \right) \mathbf{b}_{i} = \sum_{i=1}^{m} \tilde{x}_{i} \mathbf{b}_{i} = B \tilde{\mathbf{x}}.$$

Since B is non-singular and we already have $B\mathbf{x}_B = \mathbf{b}$, it follows that $\tilde{\mathbf{x}} = \mathbf{x}_B$. Thus by (2.27),

$$z \le \sum_{i=1}^m x_{B_i} c_{B_i} = z_0$$

for all x in the feasible region.

Example 2.4. Let us consider the LPP with

$$A = \begin{bmatrix} 1 & 2 & 3 & 4 \\ 1 & 0 & 0 & 1 \end{bmatrix}, \quad \mathbf{c}^T = [2, 5, 6, 8] \quad \text{and} \quad \mathbf{b} = \begin{bmatrix} 5 \\ 2 \end{bmatrix}.$$

Let us choose our starting B as

$$B = [\underbrace{\mathbf{a}_1, \mathbf{a}_4}] = [\underbrace{\mathbf{b}_1, \mathbf{b}_2}] = \begin{bmatrix} 1 & 4 \\ 1 & 1 \end{bmatrix} \ .$$

Then it is easily checked that the corresponding basic solution is $x_B = [1, 1]^T$, which is clearly feasible with objective value

$$z = \mathbf{c}_B^T \mathbf{x}_B = \begin{bmatrix} 2, 8 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = 10.$$

Since

$$B^{-1} = \frac{1}{3} \begin{bmatrix} -1 & 4 \\ 1 & -1 \end{bmatrix}, \qquad A = B \Upsilon$$

by (2.24), the y_{ij} are given by

$$\mathbf{y}_{1} = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \quad \mathbf{y}_{2} = \begin{bmatrix} -\frac{2}{3} \\ \frac{2}{3} \end{bmatrix} \quad \mathbf{y}_{3} = \begin{bmatrix} -1 \\ 1 \end{bmatrix} \quad \mathbf{y}_{4} = \begin{bmatrix} 0 \\ 1 \end{bmatrix}.$$

Hence

$$z_2 = [2, 8] \begin{bmatrix} -\frac{2}{3} \\ \frac{2}{3} \end{bmatrix} = 4,$$

and

$$z_3 = [2, 8] \begin{bmatrix} -1 \\ 1 \end{bmatrix} = 6.$$

Since $z_2 - c_2 = -1$ and $z_3 - c_3 = 0$, we see that a_2 is the *entering* column. As remarked above,

$$z_1 - c_1 = z_4 - c_4 = 0$$
,

because x_1 and x_4 are basic variables. Looking at the column entries of y_2 , we find that y_{22} is the only positive entry. Hence $b_2 = a_4$ is the *leaving* column. Thus

$$\hat{B} = [\mathbf{a}_1, \mathbf{a}_2] = \begin{bmatrix} 1 & 2 \\ 1 & 0 \end{bmatrix},$$

and the corresponding basic solution is found to be $\hat{\mathbf{x}}_B = [2, \frac{3}{2}]^T$, which is clearly feasible as expected. The new objective value is given by

$$\hat{z} = [2, 5] \begin{bmatrix} 2\\ \frac{3}{2} \end{bmatrix} = 11.5 > z.$$

$$R = \begin{cases} 1 & 4 \\ 4 & 6$$

$$B = \begin{bmatrix} 1 & 4 \\ 1 & 1 \end{bmatrix}$$

 $(\vec{a}, \vec{a}, \vec{a}, \vec{a}) = (\vec{a}, \vec{a}) (\vec{b}, +, +, \hat{a})$

CB = (C', C4) = (5 8)

$$\begin{cases}
1 & 4 \\
1 & 1
\end{cases}
\begin{pmatrix}
X_1 \\
X_4
\end{pmatrix} = \begin{pmatrix}
5 \\
2 & 1
\end{cases}$$

$$\begin{cases}
X_1 \\
X_4
\end{pmatrix} = \begin{pmatrix}
5 \\
2
\end{pmatrix}$$

$$\begin{cases}
X_1 \\
X_4
\end{pmatrix} = \begin{pmatrix}
1 & 4 \\
2
\end{pmatrix}$$

"Corresponding to the basic Variables V's Columns = I.

$$\vec{z} = (2,8) \cdot \vec{y} = (2,4,6,8)$$

$$\vec{C} - \vec{z}^{T} = (2,568) - (2+68)$$

$$= (0, 1, 0, 0)$$
 if $r_2 - z_2 > 0$
 X_2 & ele entery variable

$$\begin{cases} \chi_1 = 1 \\ \chi_1 = 0 \\ \chi_3 = 0 \\ \chi_4 = 1 \end{cases} \qquad \begin{cases} \chi_1 > 0 \\ \chi_2 > 0 \\ \chi_4 = 0 \end{cases}$$

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Since

$$\hat{B}^{-1} = \frac{1}{2} \begin{bmatrix} 0 & 2 \\ 1 & -1 \end{bmatrix},$$

by (2.24), the \hat{y}_{ij} are given by

$$\hat{\mathbf{y}}_{1} = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \quad \hat{\mathbf{y}}_{2} = \begin{bmatrix} 0 \\ 1 \end{bmatrix} \quad \hat{\mathbf{y}}_{3} = \begin{bmatrix} 0 \\ \frac{3}{2} \end{bmatrix} \quad \hat{\mathbf{y}}_{4} = \begin{bmatrix} 1 \\ \frac{3}{2} \end{bmatrix}. \qquad \qquad \hat{\mathbf{A}} = \hat{\mathbf{B}}_{1} \hat{\mathbf{Y}}_{1}$$

Hence

$$z_3 - c_3 = [2, 5] \begin{bmatrix} 0 \\ \frac{3}{2} \end{bmatrix} - 6 = 1.5,$$

and

$$z_4 - c_4 = [2, 5] \begin{bmatrix} 1 \\ \frac{3}{2} \end{bmatrix} - 8 = 1.5.$$

Since all $z_j - c_j \ge 0$, $1 \le j \le 4$, we see that the point $[2, \frac{3}{2}]$ is an optimal solution.

The last example illustrates how one can find the optimal solution by searching through the basic feasible solutions. That is exactly what *simplex method* does. However, the simplex method uses tableaus to minimize the book-keeping work that we encountered in the last example.

port for XI, XI as basic $\hat{B} = \begin{pmatrix} 1 & 2 \\ 1 & 0 \end{pmatrix}$ $\frac{1}{2} = \frac{1}{1} = \frac{1}$ $\frac{\Delta}{Z} = \frac{\hat{C}_B}{\hat{C}_B} = (2 5) \hat{Y} =$ C - Z = (2568) - (255)= (0,0,-1.5, -1.5) = "optime" None are positive => 73 & X4 V "objective à linea"

Th 2.7

 $C_{j}-2_{j} \leq 0$ $\forall j \Rightarrow tlen we are at the sphimum$

 $\vec{C}_{B}\vec{X}_{B} = \vec{C}(\vec{X}_{B}) = cmeet objective value$

 $\forall \vec{u} \in FR: \vec{C} \cdot \vec{u} = \sum_{i=1}^{n} C_i u_i \in \sum_{j=1}^{n} Z_j u_j$ (by assumption)

 $= \overline{Z}$ $= \overline{Z}$ $= \overline{Z}$ $= \overline{Z}$ $= \overline{Z}$ $= \overline{Z}$

 $= \vec{C}_{A} \vec{Y} \vec{u} \quad (A=B\vec{Y})$

= C3 BA 1 (Aa= 6)

7 CB XB

Max
$$X_0 = 3X_1 + X_2 + 3X_3 + 0X_4 + 0X_5 + 0X_6$$

Toprewhile

2 1 1 0 0 | X_1

2 2 1 0 0 | X_2

3 X_4

3 X_5

4 X_5

4 X_5

4 X_5

4 X_5

5 | Deni Variable

A | Deni Variable

Simplex

+ablean